For each $v \in V - S$, choose $u \in V - S$ such that $(u, v) \in E$. // minimum

For each $v \in V - S$, \( R[v] = R(v, v) // u = v \).

Given \( \gamma : V \rightarrow \mathbb{N}_0 \)

\( \gamma(v, v) = R[v] \), for $v \in V$.

A function $\gamma : V \rightarrow \mathbb{N}_0$. $\gamma$ is indegree of $G$.

Example: Directed graph $G = (V, E)$, $G$ source node $v$.

Simple source problem.

Shortest paths (Dijkstra's, A*).
For any inductive set $S$, prove:

\[ P \]
Suppose a vertex $v$, its
If $DC = 0$, then it is a branch.
If $DC > 0$, then $v$ is a

The shortest path from $v$ to $w

$E \in \mathcal{E}$ and $p, q, r$ are

1. a vertex

\[
\theta
c$

\pi \geq \lambda(c)

\pi \geq \lambda(c)

\pi \geq \lambda(c)

\pi \geq \lambda(c)$
\[ p_0 = (p, r, p) \]

Let \( t \in L \) be a linear character of \( S \), and let

\[ (c_1, c_2) \]


Corollary

For \( c = \lambda \),\( c = \kappa \), and for every \( \lambda, \kappa \),

\[ \text{Claim: } (\lambda \circ \kappa) = \lambda \circ \kappa \]

Proof.

By the definition of \( \lambda, \kappa \).
(c) \[ \exists \sigma \ 	ext{Lsat} \ + \ \phi \ \text{expands} \ P_{es} \]